

Lecture 11 - February 28

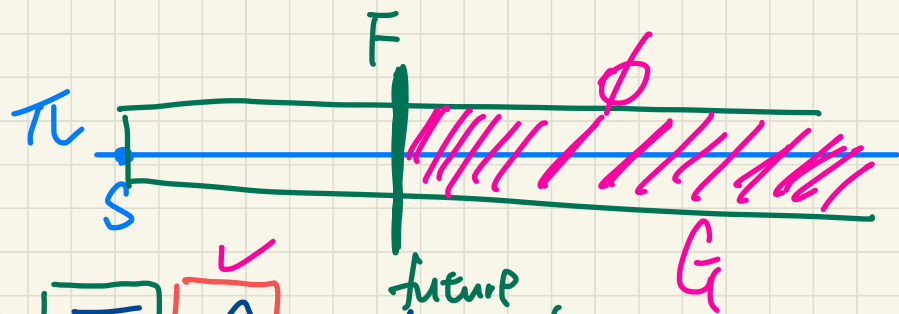
Model Checking

Path Satisfaction: Nested LTL Operators
FG vs. $F \Rightarrow FG$

Announcements

- Released: **WrittenTest1**, **Lab2** solution
- To be released:
 - + **ProgTest1** Guide (by the end of Wednesday)
 - + **ProgTest1** practice questions (by Thursday class)

- 1 ~ 2 algorithms
- conditions, loops, tuples
- assertions (postcondition)



S

=

F

G

ϕ

some state

model satisfaction
 (consider all paths starting with s).

future point

arbitrary
 LTL formula

G

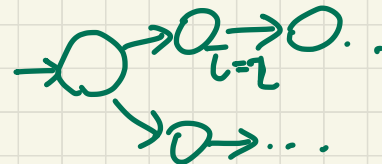
$$\forall x \cdot P(x)$$

disprove:
- find an x st. $\neg P(x)$

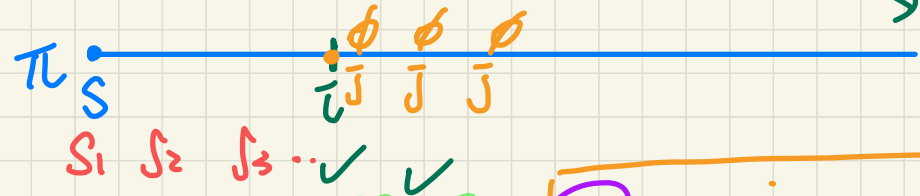
$$\exists x \cdot P(x)$$

disprove:
- find all possible x st. $\neg P(x)$

Nesting "Global" and "Future" in LTL Formulas



$$\underline{s} \models \mathbf{FG} \phi$$



Each path starting with s is s.t. eventually, ϕ holds continuously.

Q. Formulate the above nested pattern of LTL operator.

specific to the being considered

$$\forall \pi \cdot \pi = S \rightarrow \dots \Rightarrow \left(\exists \bar{i} \cdot \bar{i} \Vdash \wedge (\forall \bar{j} \cdot \bar{j} \Vdash \bar{i} \Rightarrow (\underline{\pi} \Vdash \phi)) \right)$$

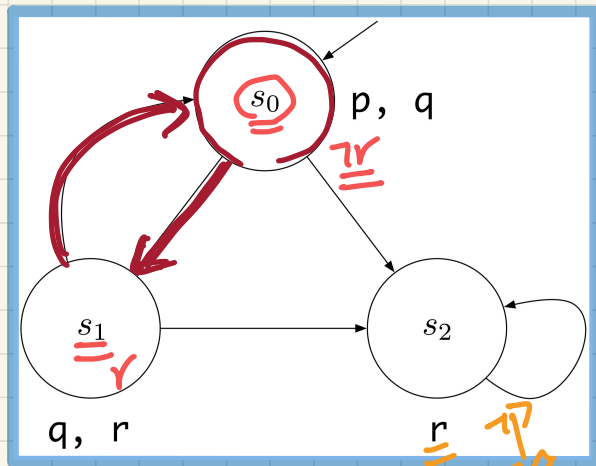
Q. How to prove the above nested pattern of LTL operators?

- * ① consider all path patterns starting from S (including i th state)
- ** ② find such \bar{i} \Rightarrow each state subsequent to \bar{i} th state satisfies ϕ

Q. How to disprove the above nested pattern of LTL operators?

- * ① Find a witness $\pi = S \rightarrow \dots$
- ** ② Show that for each state in π , \exists there's one subsequent state that violates ϕ .

Path Satisfaction: Exercises (5.1)



$s \models \phi \Leftrightarrow$ all π starting at s , $\pi \models \phi$

$s_0 \models \mathbf{FG} r$ **false**

Witness: $s_0 \rightarrow s_1 \rightarrow s_0 \rightarrow s_1 \rightarrow \dots$

$s_0 \models \mathbf{FG} (p \vee q)$ **false**

violates r , cannot be ignored in path

Witness: $s_0 \rightarrow s_2 \rightarrow s_2 \rightarrow \dots$

States involved in all possible paths satisfy prop.

$s_0 \models \mathbf{FG} (p \vee r)$ **True**

get stuck here, and both p and q are violated

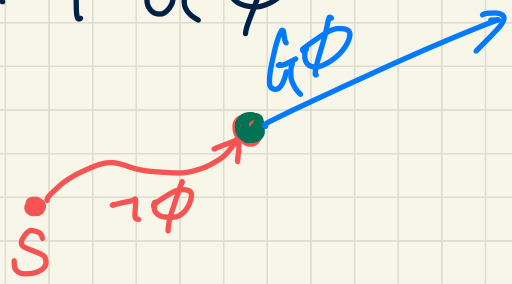
- ① $s_0 \rightarrow s_2 \rightarrow s_2 \rightarrow \dots$
- ② $s_0 \rightarrow s_1 \rightarrow s_2 \rightarrow s_2 \rightarrow \dots$
- ③ $s_0 \rightarrow s_1 \rightarrow s_0 \rightarrow s_1 \rightarrow \dots$

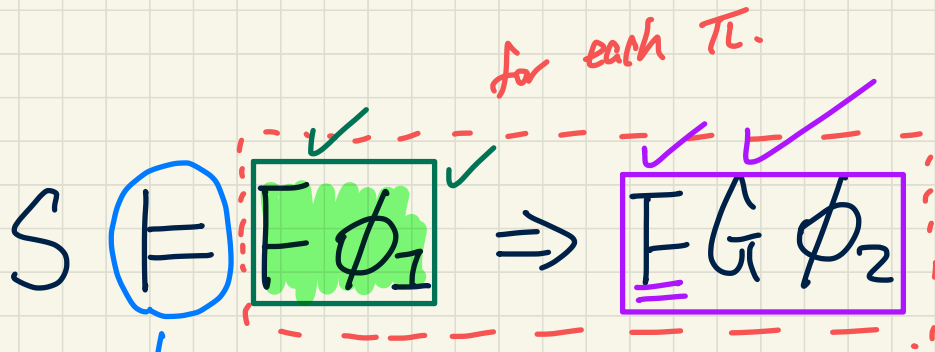
Exercise: What if we change the LHS to s_2 ?

$$S \models G\phi$$

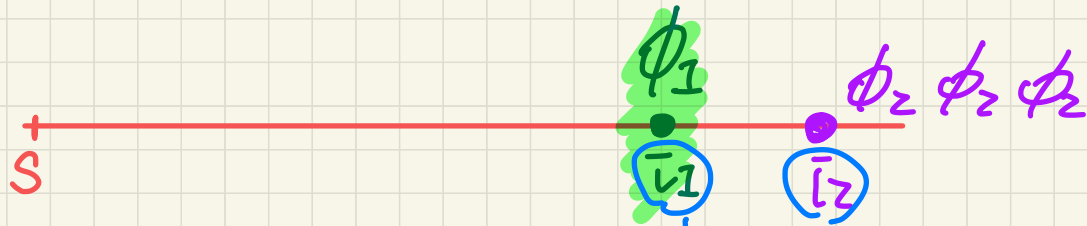


$$S \models FG\phi$$





Consider
all paths
 $\pi = S \rightarrow \dots$



- ① $\bar{t}_1 = \bar{t}_2$
- ② $\bar{t}_1 > \bar{t}_2$
- ③ $\bar{t}_1 < \bar{t}_2$

may or may
not be the same

Nesting "Global" and "Future" in LTL Formulas

$$s \models \boxed{F\phi_1} \Rightarrow \boxed{FG\phi_2}$$

Each path π starting with s is s.t. if eventually ϕ_1 holds on π , then ϕ_2 eventually holds on π continuously.

Q. Formulate the above nested pattern of LTL operators.

disproves at different scopes

$$\forall \pi \cdot \pi = s \rightarrow \dots \Rightarrow \left(\begin{aligned} &(\exists \bar{t} \cdot \bar{t} \geq 1 \wedge \pi_{\bar{t}} \models \phi_1) \\ &\Rightarrow (\exists \bar{t}_1, \bar{t}_2 \cdot \bar{t}_1 \geq 1 \wedge (\forall j \cdot j \geq \bar{t}_1 \Rightarrow \pi_j \models \phi_2)) \end{aligned} \right)$$

Q. How to **prove** the above nested pattern of LTL operators?

Q. How to **disprove** the above nested pattern of LTL operators?